

# On the 9-total-colorability of planar graphs with maximum degree 8 and without intersecting triangles

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## ABSTRACT

Recently, Kowalik, Sereni, and Škrekovski proved that planar graphs with maximum degree 9 are 10-totally colorable. This work proves that planar graphs with maximum degree 8 and without intersecting triangles are 9-totally colorable.

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## 1. Introduction

All graphs considered in this work are finite, simple and undirected. Undefined terminology and notation in this work can be found in the book by Bondy and Murty [1].

A graph is *planar* if it can be embedded into the plane so that its edges only meet at their ends. Any such concrete embedding of a planar graph is called a *plane* graph. For a plane graph  $G$ , we denote its vertex set, edge set, face set, maximum degree and minimum degree by  $V(G)$ ,  $E(G)$ ,  $F(G)$ ,  $\Delta(G)$  and  $\delta(G)$  (or simply  $V$ ,  $E$ ,  $F$ ,  $\Delta$  and  $\delta$ ), respectively. A  $k$ -cycle is a cycle of length  $k$ . A 3-cycle is also called a triangle. Two triangles are *intersecting* if they have at least one vertex in common.

A  $k$ -total-coloring of a graph  $G$  is a mapping  $\phi$  from  $V \cup E$  to the set of colors  $\{1, 2, \dots, k\}$  such that  $\phi(x) \neq \phi(y)$  for every pair of adjacent or incident elements  $x, y \in V \cup E$ . The graph  $G$  is called  $k$ -totally colorable if it has a  $k$ -total-coloring. It is clear that at least  $\Delta + 1$  colors are needed to color  $G$  totally. Vizing [2] and Behazd [3] independently conjectured that every graph is  $(\Delta + 2)$ -totally colorable. This conjecture is known as the Total Coloring Conjecture (TCC).

TCC was extensively researched in the literature, see [4–9]. Even for planar graphs, TCC remains open, see [10–13,9]. So far, we have not found any planar graph with maximum degree  $\Delta \geq 4$  which is not  $(\Delta + 1)$ -totally colorable. On the total colorability of planar graphs, one often hopes to get the best result, namely, to prove that the planar graphs under consideration are  $(\Delta + 1)$ -totally colorable. Borodin et al. [14], Wang [16] and Kowalik et al. [15] successively proved that planar graphs with  $\Delta \geq 11$ ,  $\Delta = 10$  and  $\Delta = 9$  are  $(\Delta + 1)$ -totally colorable. As far as we know, it is unknown whether all planar graphs with maximum degree 8 are 9-totally colorable! This work shows a moderate result on the topic as follows:

**Theorem 1.** Every planar graph  $G$  with  $\Delta = 8$  and without intersecting triangles is 9-totally colorable.

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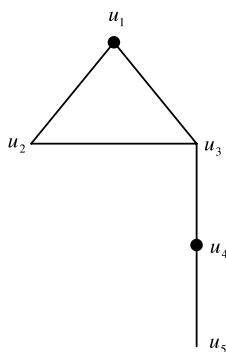


Fig. 1. Forbidden configurations in  $G$ .

## 2. Reducible configurations

Let  $G$  be a counterexample to Theorem 1 with  $\sigma(G) = |V| + |E|$  as small as possible. That is, any proper subgraph of  $G$  is 9-totally colorable while  $G$  itself is not. Embed  $G$  into the plane. It is easy to see that  $G$  is 2-connected. Hence,  $G$  has no vertices of degree 1 and the boundary of every face of  $G$  is a cycle.

For a face  $f \in F$ , its degree, denoted by  $d(f)$ , is defined to be the length of its boundary walk. Call  $f$  a  $k$ -face, a  $k^+$ -face, or a  $k^-$ -face if  $d(f) = k$ ,  $d(f) \geq k$ , or  $d(f) \leq k$ , respectively. The concepts of a  $k$ -vertex, a  $k^+$ -vertex, and a  $k^-$ -vertex are similarly defined.

**Lemma 1.** (a) Let  $uv \in E(G)$ . If  $d(u) \leq 4$ , then  $d(u) + d(v) \geq 10$ .

(b) The subgraph induced by all edges joining 2-vertices to  $\Delta$ -vertices in  $G$  is a forest.

The proof of Lemma 1 can be found in [14]. As corollaries of Lemma 1(a), the two neighbors of a 2-vertex are 8-vertices; the neighbors of a 3-vertex are  $7^+$ -vertices; and the neighbors of a 4-vertex are  $6^+$ -vertices in  $G$ .

Let  $H$  be the forest stated in Lemma 1(b) and  $T$ , a maximal tree in  $H$ . It is easy to see that all leaves of  $T$  are 8-vertices. It is easy to prove by induction that  $T$  has a maximum matching, say,  $M$ , that matches every 2-vertex of  $T$  and all 8-vertices but one that is  $M$ -unmatched. The unmatched 8-vertex by  $M$  is called the root of  $T$ . Fix a maximum matching for each maximal tree in  $H$ . Let  $v$  be a 2-vertex in  $G$ . The neighbor of  $v$  which is matched to  $v$  by the prescribed matching is called the master of  $v$ . The concept of the master was first introduced in [14] and will play an important role in the next section.

**Lemma 2.**  $G$  contains no configurations depicted in Fig. 1 where the vertices marked by  $\bullet$  have no other neighbors in  $G$ .

**Proof.** Suppose that  $G$  has the configurations depicted in Fig. 1. By the minimality of  $\sigma(G)$ ,  $G' = G - u_1u_3$  has a 9-total-coloring  $\phi : V(G') \cup E(G') \rightarrow \{1, 2, \dots, 9\}$ . Erase the colors of  $u_1$  and  $u_4$ . Let  $S(u)$  be the set of colors used by  $\phi$  for edges incident with  $u$ , and  $\bar{S}(u) = S(u) \cup \{\phi(u)\}$ . If  $\phi(u_1u_2) \in \bar{S}(u_3)$ , then there are at most eight forbidden colors for  $u_1u_3$ , i.e.,  $u_1u_3$  can be properly colored. Clearly,  $u_1$  and  $u_4$  can be properly recolored easily. Hence,  $G$  is 9-totally colorable, a contradiction. So, we can assume that  $\phi(u_1u_2) \notin \bar{S}(u_3)$ . Without loss of generality, assume  $\bar{S}(u_3) = \{1, 2, \dots, 8\}$ ,  $\phi(u_1u_2) = 9$ . If  $\phi(u_4u_5) \neq 9$ , color  $u_1u_3$  with  $\phi(u_3u_4)$  and recolor  $u_3u_4$  with 9, then  $G$  is 9-totally colorable, a contradiction. Assume  $\phi(u_4u_5) = 9$ . At this time, we first recolor edges  $u_1u_2$ ,  $u_2u_3$ ,  $u_3u_4$  with  $\phi(u_2u_3)$ , 9,  $\phi(u_2u_3)$ , respectively, then color  $u_1u_3$  with  $\phi(u_3u_4)$ . It follows that  $G$  is 9-totally colorable, a contradiction showing that  $G$  has no configuration in Fig. 1.  $\square$

Call a 3-face *bad* if it has a  $3^-$ -vertex. Similarly, a 4-face is *bad* if it has two  $3^-$ -vertices (note that they are non-adjacent by Lemma 1).

**Lemma 3.** Let  $v$  be an 8-vertex incident with a bad 3-face  $T$ . If  $v$  is adjacent to a 2-vertex not on  $T$ , then  $v$  is incident with at most five bad 4-faces.

**Proof.** Note that if  $v$  is incident with more than five bad 4-faces, then  $v$  is in one of the six configurations depicted in Fig. 2. Suppose that  $G$  has the configuration shown in Fig. 2(1). By the minimality of  $\sigma(G)$ ,  $G' = G - vv_1$  has a 9-total-coloring  $\phi : V(G') \cup E(G') \rightarrow \{1, 2, \dots, 9\}$ . Erase the colors of all the adjacent  $3^-$ -vertices  $v_1, v_2, v_3, v_4, v_5, v_6$  and  $v_8$ .

If  $\phi(v_1u_1) \in \bar{S}(v)$ , then the forbidden colors for  $vv_1$  are at most 8, i.e.,  $vv_1$  can be properly colored. After recoloring the erased vertices, we obtain a 9-total-coloring of  $G$ , a contradiction. If  $\phi(v_1u_1) \notin \bar{S}(v)$ , we can assume  $\bar{S}(v) = \{1, 2, \dots, 8\}$  with  $\phi(vv_i) = i - 1, i = 2, 3, \dots, 8, \phi(v) = 8$  and  $\phi(v_1u_1) = 9$ . Note that  $\phi(u_1v_8) \neq 9$ . If  $\phi(v_7v_8) \neq 9$ , color  $vv_1$  with  $7 = \phi(vv_8)$  and recolor  $vv_8$  with 9, getting a 9-total-coloring of  $G$ , a contradiction. So, we can assume  $\phi(v_7v_8) = 9$ . At this time,  $\phi(v_7) \neq 9$ . We first recolor  $v$  with 9 and then color  $vv_1$  with 8, easily getting a 9-total-coloring of  $G$ , a contradiction showing that  $G$  has no configuration in Fig. 2(1).

All the configurations shown in Fig. 2(2)–(6) can be similarly proved. We omit the details here.  $\square$

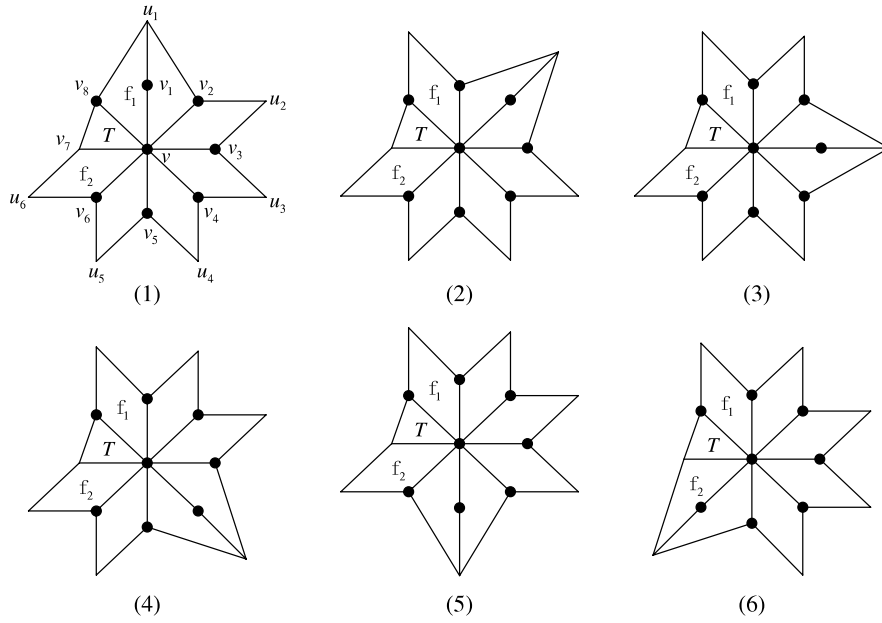


Fig. 2.  $v$  is incident with a bad 3-face and six bad 4-faces.

### 3. Discharging

To complete the proof of Theorem 1, we shall derive a contradiction by a discharging procedure based on the structural properties of  $G$  that were established in Section 2. The initial charge function  $ch$  in the discharging procedure is defined as

$$ch(x) = \begin{cases} 2d(x) - 6, & \text{if } x \in V; \\ d(x) - 6, & \text{if } x \in F. \end{cases}$$

By Handshaking Lemmas  $\sum_{v \in V} d(v) = 2|E| = \sum_{f \in F} d(f)$  and Euler's formula  $|V| - |E| + |F| = 2$ , we have

$$\sum_{x \in V \cup F} ch(x) = \sum_{x \in V} (2d(x) - 6) + \sum_{x \in F} (d(x) - 6) = -12.$$

Since any discharging procedure preserves the total charge of  $G$ , if we can define suitable discharging rules to change the initial charge function to the final charge function  $ch'$  on  $V \cup F$  such that  $ch'(x) \geq 0$  for all  $x \in V \cup F$ , then we get an obvious contradiction  $0 \leq \sum_{x \in V \cup F} ch'(x) = \sum_{x \in V \cup F} ch(x) = -12$ , which completes our proof.

Now, we define the discharging rules as follows:

R1: Charge to a 2-vertex  $v$ .

R1.1. If  $v$  is on a 3-face, then it gets 1 from each of its incident vertices.

R1.2. If  $v$  is not on a 3-face, then it gets 2 from its master.

R2: Charge to a 3-face.

R2.1. A 3-face incident with a  $4^-$ -vertex gets  $\frac{3}{2}$  from each of its two incident  $6^+$ -vertices.

R2.2. A 3-face not incident with any  $4^-$ -vertex gets 1 from each of its incident vertices.

R3: Charge to a 4-face.

R3.1. A bad 4-face gets 1 from each of its two incident  $7^+$ -vertices.

R3.2. A 4-face incident with only one  $3^-$ -vertex gets  $\frac{3}{4}$  from each of its two incident  $7^+$ -vertices;  $\frac{1}{2}$ , from the left incident  $4^+$ -vertex.

R3.3. A 4-face not incident with any  $3^-$ -vertex gets  $\frac{1}{2}$  from each of its incident vertices.

R4: Charge to a 5-face.

R4.1. A 5-face incident with two  $3^-$ -vertices gets  $\frac{1}{3}$  from each of its incident  $7^+$ -vertices.

R4.2. A 5-face incident with only one  $3^-$ -vertex gets  $\frac{1}{4}$  from each of its incident  $4^+$ -vertices.

R4.3. A 5-face not incident with any  $3^-$ -vertex gets  $\frac{1}{5}$  from each of its incident vertices.

These rules are illustrated in Fig. 3.

The rest of this work is devoted to checking that  $ch'(x) \geq 0$  for all  $x \in V \cup F$ . First note that our discharging rules are just designed so that  $ch'(f) \geq 0$  for all  $f \in F$  and  $ch'(v) \geq 0$  for all 2-vertices in  $V$ . So we only need to check that  $ch'(v) \geq 0$  for all  $3^+$ -vertices in  $G$ .

**Remark 1.**  $\forall v \in V$ , there is at most one 3-face incident with  $v$  for  $G$  having no intersecting triangles.

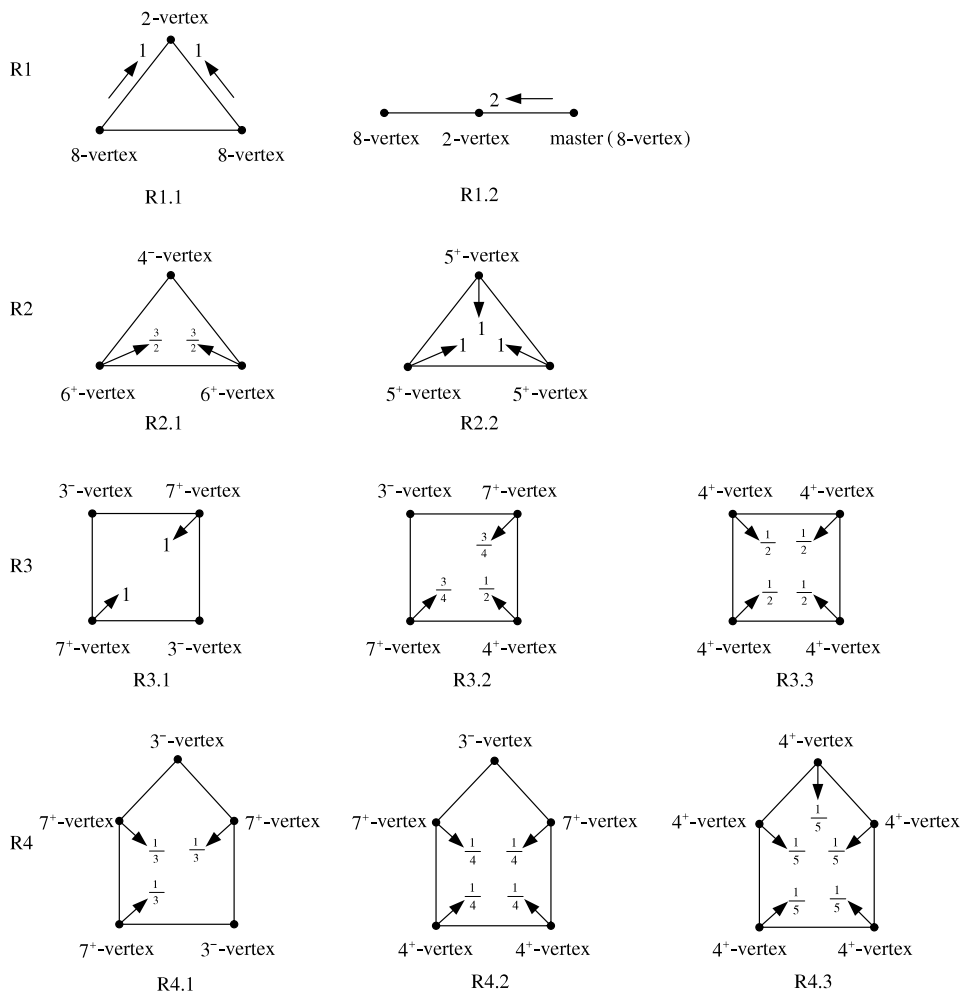


Fig. 3. Discharging rules.

**Remark 2.** A  $7^+$ -vertex discharges nothing to each of its incident  $6^+$ -faces; at most 1 to each of its incident bad 4-faces; at most  $\frac{3}{4}$  to each of its incident non-bad 4-faces.

Let  $d(v) = 3$ . According to our discharging rules, no charge is discharged to or from  $v$ , that is,  $ch'(v) = ch(v) = 0$ .

Let  $d(v) = 4$ . By R2.1, R3.2, R3.3, R4.2 and R4.3,  $v$  discharges at most  $\frac{1}{2}$  to each of its incident faces. So,  $ch'(v) \geq ch(v) - \frac{1}{2} \times 4 = 0$ .

Let  $d(v) = 5$ . By our rules,  $v$  discharges at most 1 to a possible incident 3-face, and at most  $\frac{1}{2}$  to each of the left  $4^+$ -faces. Thus,  $ch'(v) \geq ch(v) - 1 - \frac{1}{2} \times 4 = 1$ .

Let  $d(v) = 6$ . By our rules,  $v$  discharges at most  $\frac{3}{2}$  to a possible incident 3-face, and at most  $\frac{1}{2}$  to each of the left  $4^+$ -faces. It follows that  $ch'(v) \geq ch(v) - \frac{3}{2} - \frac{1}{2} \times 5 = 2$ .

Let  $d(v) = 7$ . By Remarks 1 and 2,  $ch'(v) \geq ch(v) - \frac{3}{2} - 1 \times 6 = \frac{1}{2}$ .

Let  $d(v) = 8$ . If  $v$  is not adjacent to any 2-vertex, then  $ch'(v) \geq ch(v) - \frac{3}{2} - 1 \times 7 = \frac{3}{2}$ . If  $v$  is not incident with any 3-faces, by R1 and Remark 2,  $ch'(v) \geq ch(v) - 2 - 1 \times 8 = 0$ . Assume that  $v$  is incident with a 3-face, say,  $T$ , and adjacent to at least one 2-vertex. If  $T$  has a 2-vertex, say,  $u$ , then, by Lemma 2,  $v$  is not incident with any other 2-vertex, namely, the maximal tree containing  $u$  in  $H$  is a path of length 2. According to R1.1,  $ch'(v) \geq ch(v) - 1 - \frac{3}{2} - 1 \times 7 = \frac{1}{2}$ .

Now, we can assume that  $T$  has no 2-vertex. If  $T$  is not a bad 3-face, then both the two faces adjacent to  $T$  and incident with  $v$ , say,  $f_1$  and  $f_2$ , are not bad 4-faces. By our rules and, in particular, by Remark 2,  $v$  discharges at most  $\frac{3}{4}$  to each of  $f_1$  and  $f_2$ . Thus,  $ch'(v) \geq ch(v) - 2 - \frac{3}{2} - 1 \times 5 - \frac{3}{4} \times 2 = 0$ . If  $T = vxy$  is a bad 3-face with a 3-vertex, say,  $x$ , by Lemma 3, all the faces incident with  $v$  other than  $T$  are  $4^+$ -faces and the number of the bad 4-faces is at most 5. By Remarks 1 and 2,  $ch'(v) \geq ch(v) - 2 - \frac{3}{2} - 1 \times 5 - \frac{3}{4} \times 2 = 0$ . This concludes the proof of the theorem.  $\square$

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